

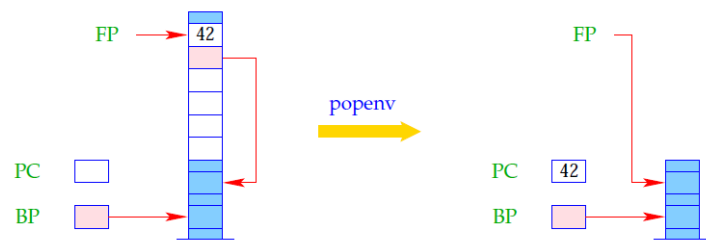
Title: Seidl: Virtual\_Machines (10.07.2012)

Date: Tue Jul 10 15:19:33 CEST 2012

Duration: 12:22 min

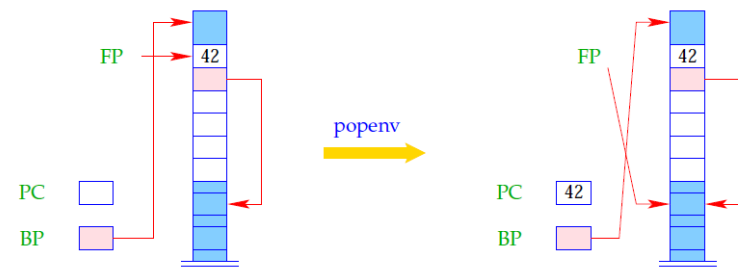
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The instruction `popenv` restores the registers `FP` and `PC` and possibly pops the stack frame:



if ( $FP > BP$ )  $SP = FP - 6$ ;  
 $PC = posCont$ ;  
 $FP = FPold$ ;

Warning: `popenv` may fail to de-allocate the frame !!!



if ( $FP > BP$ )  $SP = FP - 6$ ;  
 $PC = posCont$ ;  
 $FP = FPold$ ;

If popping the stack frame fails, new data are allocated on top of the stack. When returning to the frame, the locals still can be accessed through the `FP` :-))

### 33 Queries and Programs

The translation of a program:  $p \equiv rr_1 \dots rr_h ? g$  consists of:

- an instruction `no` for failure;
- code for evaluating the query  $g$ ;
- code for the predicate definitions  $rr_i$ .

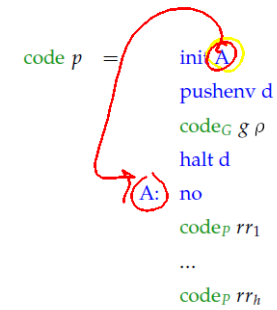
Preceding query evaluation:

- ⇒ initialization of registers
- ⇒ allocation of space for the globals

Succeeding query evaluation:

- ⇒ returning the values of globals

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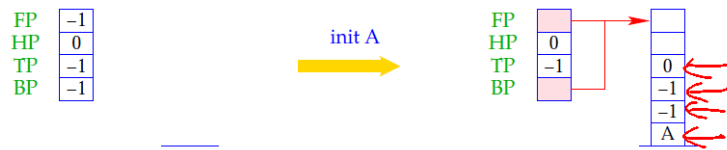
where  $free(g) = \{X_1, \dots, X_d\}$  and  $\rho$  is given by  $\rho X_i = i$ .

The instruction `halt d` ...

- ... terminates the program execution;
- ... returns the bindings of the  $d$  globals;
- ... causes backtracking — if demanded by the user :-)

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The instruction `init A` is defined by:

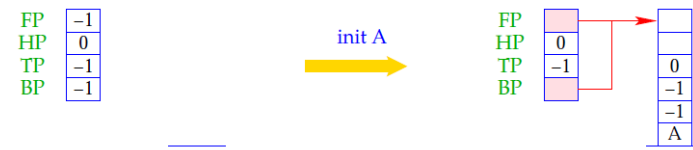


BP = FP = SP = 5;  
 S[0] = A;  
 S[1] = S[2] = -1;  
 S[3] = 0;  
 BP = FP;

At address "A" for a failing goal we have placed the instruction `no` for printing `no` to the standard output and `halt :-)`

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The instruction `init A` is defined by:



BP = FP = SP = 5;  
 S[0] = A;  
 S[1] = S[2] = -1;  
 S[3] = 0;  
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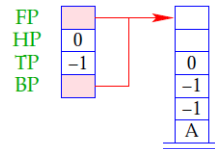
At address "A" for a failing goal we have placed the instruction `no` for printing `no` to the standard output and `halt :-)`

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The instruction `init A` is defined by:

FP	-1
HP	0
TP	-1
BP	-1

`init A`



BP = FP = SP = 5;  
 S[0] = A;  
 S[1] = S[2] = -1;  
 S[3] = 0;  
 BP = FP;

At address "A" for a failing goal we have placed the instruction `no` for printing `no` to the standard output and `halt :-)`

### The Final Example:

$t(X) \leftarrow \bar{X} = b$        $q(X) \leftarrow s(\bar{X})$        $s(X) \leftarrow \bar{X} = a$   
 $p \leftarrow q(\bar{X}), t(\bar{X})$        $s(X) \leftarrow t(\bar{X})$       ?  $p$

The translation yields:

<code>init N</code>	<code>popenv</code>	<code>q/1: pushenv 1</code>	<code>E: pushenv 1</code>
<code>pushenv 0</code>	<code>p/0: pushenv 1</code>	<code>mark D</code>	<code>mark G</code>
<code>mark A</code>	<code>mark B</code>	<code>putref 1</code>	<code>putref 1</code>
<code>call p/0</code>	<code>putvar 1</code>	<code>call s/1</code>	<code>call t/1</code>
<code>A: halt 0</code>	<code>call q/1</code>	<code>D: popenv</code>	<code>G: popenv</code>
<code>N: no</code>	<code>B: mark C</code>	<code>s/1: setbtp</code>	<code>F: pushenv 1</code>
<code>t/1: pushenv 1</code>	<code>putref 1</code>	<code>try E</code>	<code>putref 1</code>
<code>putref 1</code>	<code>call t/1</code>	<code>delbtp</code>	<code>uatom a</code>
<code>uatom b</code>	<code>C: popenv</code>	<code>jump F</code>	<code>popenv</code>